

## FPGA Implementation of 16 bit MUX Based Multiplier

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### Abstract

In digital filter implementation, the multiplier usage is avoided by using MUX based multiplier and Look up Table (LUT) based multiplier. Odd multiple storage proposed by stores the product of odd multiple of co-efficient and the input. The advantage of storing odd multiple is that even multiples can be obtained by a simple left shift operation. 8 bit MUX based multiplication is carried out only with 16 bit (2x1) MUX and shifters. Proposed 16 bit MUX based multiplication is carried out only with 32 bit (2x1) MUX and shifters. The designs've been implemented using Verilog and synthesized using Xilinx 13.2 Spartan 3e kit.

**Keywords:** *FIR filter, Look-up Table, Reconfigurable Architecture, and Distributed Arithmetic.*

### 1. Introduction

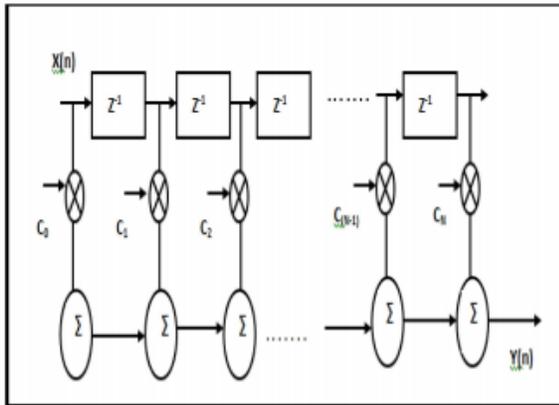
Finite Impulse Response (FIR) digital filter is widely used as a basic block in signal and image processing applications. The number of multiply-accumulate (MAC) operations required per filter output increases linearly with the order of filter, but implementation of higher order filters in real-time is another challenging task. Recently with the advent of software defined radio (SDR), the research has been concentrated on realization of FIR filter [1] mainly due to its high flexibility and low complexity [5] [6] [7]. The digit-based reconfigurable architecture presented in [3] provides a flexible and low power solution with a wide range of precision and tap length of FIR filters. Conventionally, the FIR filters are designed based on programmable multiply-accumulate [MAC] architecture [6] and systolic architecture [7]. The performances of the designs are analyzed in terms of hardware complexity, power consumption and throughput. In [6], the programmable MAC architectures consume low power with reduced supply voltage and it requires large area. In [7], even though systolic based architecture reduces the complexity, it increases the latency when the order of the filter increases. Several attempts have been made and it continued to develop low-complexity dedicated VLSI systems for these filters.

There are several issues in the hardware implementation of Digital filters. The direct implementation of N-tap FIR filters which requires N MAC operations are too expensive to implement in hardware due to its logic complexity and area requirement. Distributed Arithmetic (DA) and LUT constitute memory-based techniques. Among them, DA is applied for inner product computation [8]-[17], LUT provides computation of multiplication [18]-[24]. In the LUT based approach, multiplications of input values with a fixed coefficient and consisting of all possible pre-computed product values corresponding to all possible values of input multiplicand, while in the DA-based approach, an LUT is used to store all possible values of inner-products of a fixed-point vector. If the inner-products are implemented directly, the memory-size of LUT-multiplier based implementation increases exponentially with the input word length, whereas the memory size of the DA-based approach increases exponentially with the inner-product-length. It is observed that the memory-size reduction is achieved by such decompositions is accompanied by raise in latency as well as the number of adders and latches.

### II. MULTIPLIER LESS STRUCTURES BASED FOR FIR FILTER

In this paper FPGA realization of the two FIR filter architectures such as Odd multiple storage scheme

[4] and MUX based multiplier are discussed. Both Odd multiple storage scheme and MUX based multiplier have multiplier less architecture to reduce complexity in the design, by manipulating the odd multiples of the fixed coefficient in LUT design, and general multipliers replaced by shift and add operations respectively. The performances of the proposed architectures are analyzed in terms of area and time complexities by varying the number of taps.



**Fig.1. General Structure of an FIR filter**

Figure 1 shows an implementation of N tap FIR filter with the usage of three elements namely adders, multipliers and delay elements. Let X (n) and Y (n) are the input and output sequences of the FIR filter respectively. Consider an N-tap FIR filter that can be formulated as

$$Y(n) = \sum_{k=0}^{N-1} h_k X(n-k) \quad (1)$$

Where  $h_k$  is the  $k$ th coefficient of the filter impulse response.

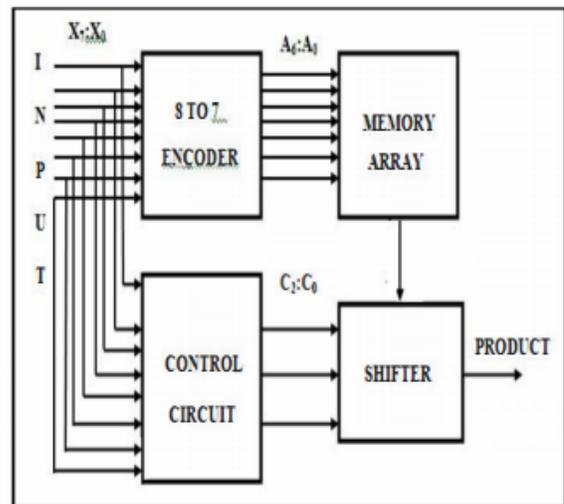
In the above equation, the filter coefficients do not vary with regard to the input signal or the noise. This implies that the coefficients are fixed. Such filters whose coefficients are fixed are called as fixed filters or filters with fixed characteristics. In the filter the MAC structure and delay blocks are the main building blocks. The performance of the filter can be enhanced by the introduction of two architecture schemes which reduces the complexity and critical path. The two schemes discussed here are ODD and MUX based multiplier.

**III. ODD MULTIPLE STORAGE SCHEME**

This method proposed by [4] stores the product of odd multiple of co-efficient and the input. The

advantage of storing odd multiple is that even multiples can be obtained by a simple left shift operation. For any input vector of N bits the number of locations in the LUT is identified by  $2^N$ . However, the odd multiple storage scheme reduces the number of locations by a factor of 2. Considering an input vector of 8 bits, uses only 128(2<sup>7</sup>) unlike 256 locations of conventional LUT. Table I gives the complete insight about the look up table of odd multiple storage scheme. The 7-bit address (A6-A0) given in Table I identifies the stored values of LUT and it also gives a complete detail about the number of shifts to be performed for the corresponding input bits Hence for the input bits '10' the LUT stored value A is shifted one time to get 2A. The control bits  $c_2, c_1, c_0$  decides the number of shifts to be performed. The Figure 2 gives the proposed structure of the odd multiple storage schemes.

The architecture has an 8 to 7 bit encoder which maps the 8-bit input word ( $x_7: x_0$ ) to a 7 bit address (A6: A0) according to the logical relations. It chooses one output address value depending upon the output obtained from the encoder. The value obtained from the memory array is an odd multiple value of the coefficient. In order to obtain the even multiple values, the shift operation is to be performed. The control unit generates the control bits according to the relations 3a to 3c



**Fig.2. odd multiple storage scheme**

Sym bol	Address A <sub>6</sub> A <sub>5</sub> A <sub>4</sub> A <sub>3</sub> A <sub>2</sub> A <sub>1</sub> A <sub>0</sub>	Stored value	Input X <sub>7</sub> X <sub>6</sub> X <sub>5</sub> X <sub>4</sub> X <sub>3</sub> X <sub>2</sub> X <sub>1</sub> X <sub>0</sub>	Product value	# of shifts	Control C <sub>2</sub> C <sub>1</sub> C <sub>0</sub>
P <sub>0</sub>	0	A	1	A	0	0
			10	2A	1	1
			100	4A	2	10
			1000	8A	3	11
			10000	16A	4	100
			100000	32A	5	101
			1000000	64A	6	110
10000000	128A	7	111			
P <sub>1</sub>	1	3A	11	3A	0	0
			110	6A	1	1
			1100	12A	2	10
			11000	24A	3	11
			110000	48A	4	100
			1100000	96A	5	101
11000000	192A	6	110			
.	.	.	.	.	.	.
P <sub>121</sub>	1111001	243A	11110011	243A	0	0
P <sub>122</sub>	1111010	245A	11110101	245A	0	0
P <sub>123</sub>	1111011	247A	11110111	247A	0	0
P <sub>124</sub>	1111100	249A	11111001	249A	0	0
P <sub>125</sub>	1111101	251A	11111011	251A	0	0
P <sub>126</sub>	1111110	253A	11111101	253A	0	0
P <sub>127</sub>	1111111	255A	11111111	255A	0	0

Table.1. Odd multiple storage scheme

IV. PROPOSED MUX BASED MULTIPLIER

Figure 3 shows a shift and add multiplier implemented with the usage of namely adders, and MUX. Since the general multipliers are replaced by MUX, shifters and adders, referred to as multiplier less implementation. Consider an eight bit multiplication of signed numbers. Here the multiplication is carried out only with MUX and shifters. For example consider two signed 8-bit numbers x and c bits, represented in two's complement form.

Table II shows the binary input X is given in powers of 2. As in equation (8) the number C is left shifted based on the given input bits. If MSB bit of input is 1, the shifted value is twos complemented. These shifted values are given through MUX and the corresponding

outputs are selected based on the input bits. Finally MUX outputs are added to obtain the product.

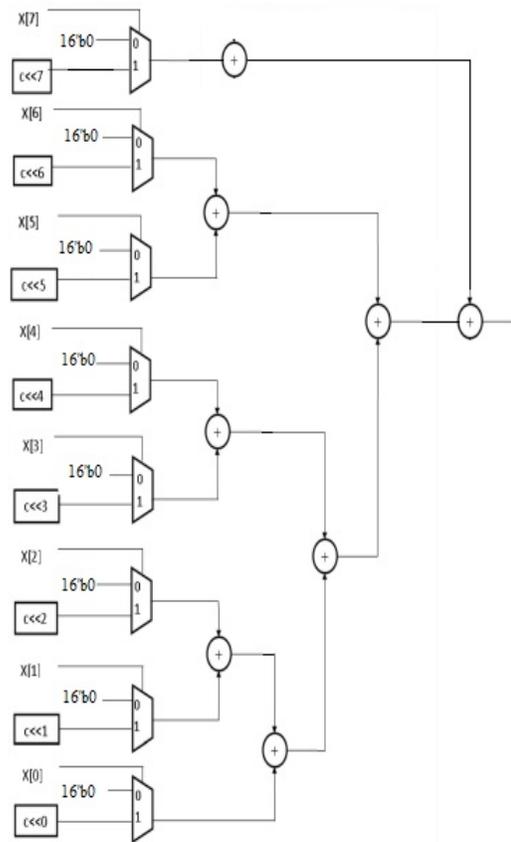


Fig.3. 8-bit MUX based Multiplier

Input, X	P = X.C	X in powers of 2
0001	1C	2 <sup>0</sup>
0010	2C	2 <sup>1</sup>
0011	3C	2 <sup>1</sup> +2 <sup>0</sup>
0100	4C	2 <sup>2</sup>
0101	5C	2 <sup>2</sup> +2 <sup>0</sup>
0110	6C	2 <sup>2</sup> +2 <sup>1</sup>
0111	7C	2 <sup>2</sup> +2 <sup>1</sup> +2 <sup>0</sup>
1000	-8C	-2 <sup>3</sup>
1001	-7C	-2 <sup>3</sup> +2 <sup>0</sup>
1010	-6C	-2 <sup>3</sup> +2 <sup>1</sup>
1011	-5C	-2 <sup>3</sup> +2 <sup>1</sup> +2 <sup>0</sup>
1100	-4C	-2 <sup>3</sup> +2 <sup>2</sup>
1101	-3C	-2 <sup>3</sup> +2 <sup>2</sup> +2 <sup>0</sup>
1110	-2C	-2 <sup>3</sup> +2 <sup>2</sup> +2 <sup>1</sup>
1111	-1C	-2 <sup>3</sup> +2 <sup>2</sup> +2 <sup>1</sup> +2 <sup>0</sup>

Table.2.MUX based multiplier

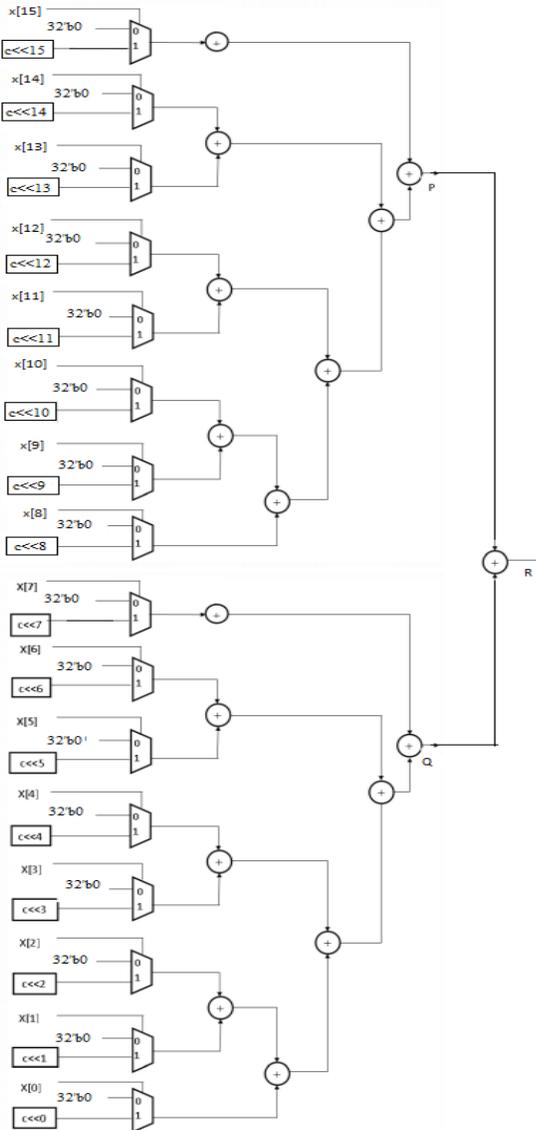


Fig.4. 16-bit MUX based Multiplier architecture

**V. CONCLUSION**

In this paper an efficient Low complexity based FIR filter architecture using MUX based multiplier and odd multiple scheme have been discussed and implemented effectively. Odd multiple storage proposed by stores the product of odd multiple of coefficient and the input. The advantage of storing odd multiple is that even multiples can be obtained by a simple left shift operation. 8 bit MUX based multiplication is carried out only with 16 bit (2x1) MUX and shifters. Proposed 16 bit MUX based multiplication is carried out only with 32 bit (2x1) MUX and shifters. The designs've been implemented

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**VI. REFERENCES**

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